

## 8th Exercise sheet for Advanced Algorithmics, SS 15

**Hand In:** Until Wednesday, 17.06.2015, 12:00am, in lecture, exercise sessions, hand-in box in stairwell 48-6 or via email.

### Problem 21

Let  $S$  a set of keys with  $|S| = n$ . Show that the expected height of a treap for  $S$  is logarithmic, i. e.

$$\mathbb{E}[\text{height}(n)] \in \mathcal{O}(\log n) .$$

**Hint:** You can approach this as follows.

1. Show that the  $\text{depth}(x) \in \mathcal{O}(\log n)$  with high probability, for all elements  $x \in S$ .
2. Show that  $\text{height}(n) \in \mathcal{O}(\log n)$  with high probability.
3. Derive the claim.

You will have to do some literature research. To get you started, here are some pointers.

- The value of Game A is also called (the expected number of) “left-to-right maxima” (in random permutations). Under this name it has been analysed deeply, e. g. in *The Art of Computer Programming Vol. 1* by Donald Knuth.
- The phrase “with high probability” has no clear definition. It often refers to bounds on the *tails* of a distribution. Such can be obtained by Markov or Chebyshev inequalities, others are called Chernoff bounds. The latter are the strongest; you can find details for instance in *Elements of Information Theory* by T. M. Cover and J. A. Thomas.
- You can think about the connection between treaps, regular binary search trees and Quicksort with respect to this analysis.

**Problem 22**

Show that the following statements hold for treaps of size  $n$ :

- a) The expected runtime of operations INSERT and DELETE is  $\mathcal{O}(\log n)$ , respectively.
- b) The expected number of rotations during an INSERT or DELETE operation is bounded above by two.

**Problem 23**

Devise an algorithm that splits treaps  $S$  as efficiently as possible into two treaps  $S_{\leq k}$  and  $S_{>k}$  which contain all keys in  $S$  that are  $\leq k$  resp.  $> k$ .

Argue why your algorithm is correct and what runtime it has.